

Infinite words containing squares at every position

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The Main Problem

We consider the following question of Richomme:

What is the infimum of the real numbers $\alpha > 2$ for which there exists an infinite word that contains no α -power as a subword, yet contains arbitrarily large squares beginning at every position?

As we shall see, over the binary alphabet, the answer to Richomme's question is $\alpha = 7/3$.

Definitions

First we recall some basic definitions.

square: A word of the form xx (like **tintin**).

cube: A word of the form xxx .

α -power: For a real number α , the shortest prefix of length $\geq \alpha|x|$ of some infinite word $xxx\cdots$.

α^+ -power: For a real number α , the shortest prefix of length $> \alpha|x|$ of some infinite word $xxx\cdots$.

overlap: A 2^+ -power.

α -power-free word: A word containing no α -power as a subword.

α^+ -power-free word: A word containing no α^+ -power as a subword.

Aperiodic Words with Arbitrarily Large Squares

- There exist aperiodic infinite binary words that contain arbitrarily large squares starting at every position.
- For instance, all Sturmian words have this property (see, for example, Allouche, Davison, Queffélec, and Zamboni 2001).
- Certain Sturmian words additionally avoid α -powers for some real number α .
- The Fibonacci word

$$\mathbf{f} = 010010100100101001010 \dots$$

contains no $(2 + \varphi)$ -powers, where $\varphi = (1 + \sqrt{5})/2$ (Mignosi and Pirillo 1992; see also Krieger 2007).

Aperiodic Words with Arbitrarily Large Squares

- However, \mathbf{f} contains arbitrarily large squares at each position.
- For example, \mathbf{f} begins with the squares
 - ▶ 010010,
 - ▶ 0100101001,
 - ▶ 0100101001001010, etc.
- By contrast, the Thue–Morse word

$$\mathbf{t} = 011010011001011010010110 \dots$$

is overlap-free (Thue 1912) but does not contain squares beginning at every position.

The Squares in the Thue–Morse Word

- Let μ denote the Thue–Morse morphism, which maps

$$0 \rightarrow 01, \quad 1 \rightarrow 10.$$

- Let $A = \{00, 11, 010010, 101101\}$ and

$$\begin{aligned} \mathcal{A} &= \bigcup_{k \geq 0} \mu^k(A) \\ &= \{00, 11, 0101, 1010, 010010, 101101, 01100110, \dots\}. \end{aligned}$$

- Then \mathcal{A} is the set of squares appearing in the Thue–Morse word (Pansiot 1981; Brlek 1989).
- The overlap-free binary squares in general are the conjugates of the words in \mathcal{A} (Thue 1912; Shelton and Soni 1985).
- (A **conjugate** of x is a word y such that $x = uv$ and $y = vu$ for some u, v .)

The Squares in Overlap-free Words

Using this characterization of the overlap-free binary squares, we generalize the previous observation that the Thue–Morse word does not contain squares starting at every position.

Theorem

If w is an infinite overlap-free binary word, then there is a position i such that w does not contain a square beginning at position i .

The Squares in Overlap-free Words

- Check by computer that any overlap-free word of length greater than 36 contains 010011.
- E.g., in the Thue–Morse word:

0110100110010110...

- Let i denote any position at which 010011 occurs in \mathbf{w} .
- We claim that no square begins at position i .
- Suppose to the contrary that xx is such a square.
- Except for 00, 11, 010010, 101101 and their conjugates, every overlap-free square xx has $|x|$ even.
- It follows that $xx = 0y10y1$, where $y \in \{01, 10\}^\ell$ for some ℓ .
- This forces xx to be followed by 0 in \mathbf{w} , so that the overlap $xx0$ occurs in \mathbf{w} , a contradiction.

A Word with Large Squares at Every Position

Theorem

There exists an infinite $(7/3)^+$ -power-free binary word that contains arbitrarily large squares beginning at every position.

- We illustrate the idea of the proof with the following example, due to Richomme.
- Let f be the map $0 \rightarrow 001, 1 \rightarrow 011$.
- f maps cubefree words to cubefree words: iterating f gives the infinite cubefree word

$$\mathbf{x} = f^\omega(0) = 001001011001001011001011011 \dots$$

The Construction Continued

- x starts with 0010010.
- x therefore starts with $f^n(001)f^n(001)f^n(0)$, for $n \geq 0$.
- $f^n(0)$ is a prefix of $f^n(001)$.
- Every subword of length $|f^n(001001)|$ starting at positions 1 to $|f^n(0)|$ is thus a square.
- For instance, when $n = 1$:

001001011001001011001011011...

- Thus x is cubefree, but contains arbitrarily large squares starting at every position.

The Construction Continued

- x starts with 0010010.
- x therefore starts with $f^n(001)f^n(001)f^n(0)$, for $n \geq 0$.
- $f^n(0)$ is a prefix of $f^n(001)$.
- Every subword of length $|f^n(001001)|$ starting at positions 1 to $|f^n(0)|$ is thus a square.
- For instance, when $n = 1$:

001001011 001001011 001011011 \dots

- Thus x is cubefree, but contains arbitrarily large squares starting at every position.

The Construction Continued

- x starts with 0010010.
- x therefore starts with $f^n(001)f^n(001)f^n(0)$, for $n \geq 0$.
- $f^n(0)$ is a prefix of $f^n(001)$.
- Every subword of length $|f^n(001001)|$ starting at positions 1 to $|f^n(0)|$ is thus a square.
- For instance, when $n = 1$:

0 0100101110 0100101110 01011011 \dots .

- Thus x is cubefree, but contains arbitrarily large squares starting at every position.

The Construction Continued

- x starts with 0010010.
- x therefore starts with $f^n(001)f^n(001)f^n(0)$, for $n \geq 0$.
- $f^n(0)$ is a prefix of $f^n(001)$.
- Every subword of length $|f^n(001001)|$ starting at positions 1 to $|f^n(0)|$ is thus a square.
- For instance, when $n = 1$:

00 100101100 100101100 1011011 \dots .

- Thus x is cubefree, but contains arbitrarily large squares starting at every position.

The Construction Continued

- x starts with 0010010.
- x therefore starts with $f^n(001)f^n(001)f^n(0)$, for $n \geq 0$.
- $f^n(0)$ is a prefix of $f^n(001)$.
- Every subword of length $|f^n(001001)|$ starting at positions 1 to $|f^n(0)|$ is thus a square.
- For instance, when $n = 1$:

001 001011001 001011001 011011 \dots .

- Thus x is cubefree, but contains arbitrarily large squares starting at every position.

The Desired Result

- To improve from cubefree to $(7/3)^+$ -power-free, we instead iterate g , which maps

0 \rightarrow 011010011011001101001

1 \rightarrow 100101100110110010110.

- Then g generates an infinite $(7/3)^+$ -power-free word (Kolpakov, Kucherov, and Tarannikov 1999) , but contains 0110110.
- As previously, iterates of g on the overlap 0110110 create arbitrarily large squares at every position.
- The details are a little more involved, but conceptually similar to the previous argument.

Optimality of the Previous Result

Theorem

If w is an infinite $7/3$ -power-free binary word, then there is a position i such that w does not contain arbitrarily large squares beginning at position i .

- The idea is to find an occurrence of 010011 at position i in w and argue that there cannot exist arbitrarily large squares at position i .
- The proof requires a characterization of the $7/3$ -power-free binary squares.
- As we shall show later, the $7/3$ -power-free binary squares are exactly the overlap-free binary squares.

Words with Squares Starting at Every Position

- What if we remove the requirement that our words contain *arbitrarily large* squares at every position?
- What if we simply want a square at every position?

Theorem

There exists an infinite $7/3$ -power-free binary word that contains squares beginning at every position.

- Define the sequence: $A_0 = 00$ and $A_{n+1} = 0\mu^2(A_n)$, $n \geq 0$.
- The first few terms are

$$A_0 = 00$$

$$A_1 = 001100110$$

$$A_2 = 0011001101001100101100110100110010110$$

⋮

Words with Squares Starting at Every Position

- As $n \rightarrow \infty$, A_n tends to an infinite limit word

$$\mathbf{w} = 0011001101001100101100110100110010110 \cdots$$

- Note: \mathbf{w} contains the squares
 - ▶ 00 at position 1,
 - ▶ 01100110 at position 2,
 - ▶ 1100110100110010 1100110100110010 at position 3, etc.
- In general, \mathbf{w} contains a square starting at every position.
- A similar, but slightly more complicated, construction yields the stronger result:

Theorem

For every real number $\alpha > 2$, there exists an infinite α -power-free binary word that contains squares beginning at every position.

Characterizing the $7/3$ -power-free Squares

- One of our main results generalizes Shelton and Soni's characterization of the overlap-free binary squares.
- Recall the sets $A = \{00, 11, 010010, 101101\}$ and

$$\begin{aligned}\mathcal{A} &= \bigcup_{k \geq 0} \mu^k(A) \\ &= \{00, 11, 0101, 1010, 010010, 101101, 01100110, \dots\}.\end{aligned}$$

Theorem

The $7/3$ -power-free binary squares are the conjugates of the words in \mathcal{A} .

- In other words, the $7/3$ -power-free binary squares are exactly the overlap-free binary squares.

The Factorization Theorem

- Recall the morphism μ , which maps $0 \rightarrow 01$ and $1 \rightarrow 10$.
- We need the factorization theorem:

Theorem (Karhumäki and Shallit 2004)

Let $x \in \{0, 1\}^$ be α -power-free, $2 < \alpha \leq 7/3$. Then there exist $u, v \in \{\epsilon, 0, 1, 00, 11\}$ and an α -power-free $y \in \{0, 1\}^*$ such that $x = u\mu(y)v$.*

- For example, if $x = 0010011001$, then $x = 00\mu(1010)$.
- Here, $u = 00$, $y = 1010$, and $v = \epsilon$.

Some Lemmas

Lemma

Let $xx \in \{0, 1\}^$ be 7/3-power-free. If $xx = \mu(y)$, then $|y|$ is even. Consequently, y is a square.*

- We also need a special version of the factorization theorem just for squares:

Lemma

Let $xx \in \{0, 1\}^$ be 7/3-power-free. If $|xx| > 8$, then either*

- (a) $xx = \mu(y)$, where $y \in \{0, 1\}^*$; or
- (b) $xx = \bar{a}\mu(y)a$, where $a \in \{0, 1\}$ and $y \in \{0, 1\}^*$.

- For example, if $xx = 010110010110$, then $xx = \mu(001001)$.
- And if, $xx = 11001100$, then $xx = 1\mu(101)0$.

A Sketch of the Proof

- Let xx be a minimal $7/3$ -power-free square that is not a conjugate of a word in \mathcal{A} .
- A computer check verifies that $|xx| > 8$.
- Case 1: $xx = \mu(y)$. Then y is a square.
- Furthermore, y is not a conjugate of a word in \mathcal{A} , contradicting the minimality of xx .
- Case 2: $xx = \bar{a}\mu(y)a$. Then $a\bar{a}\mu(y) = \mu(ay)$ is also a square zz .
- We proceed by showing that zz is $7/3$ -power-free, and consequently, that ay is a $7/3$ -power-free square, contradicting the minimality of xx .
- The details of showing that zz is $7/3$ -power-free are somewhat technical, and we omit them.

Summary of Results

- An infinite overlap-free binary word cannot contain squares beginning at every position.
- There exists an infinite $(7/3)^+$ -power-free binary word that contains arbitrarily large squares beginning at every position.
- An infinite $7/3$ -power-free binary word cannot contain arbitrarily large squares beginning at every position.
- For every real number $\alpha > 2$, there exists an infinite α -power-free binary word that contains squares beginning at every position.
- The $7/3$ -power-free binary squares are exactly the overlap-free binary squares.

Open Problems

- We have only considered words over a binary alphabet. Do similar results hold over a larger alphabet?
- For instance, does there exist an infinite overlap-free ternary word that contains squares beginning at every position?
- Richomme observed that over any alphabet there cannot exist an infinite overlap-free word containing infinitely many squares at every position.
- Apply the following result of Ilie (2007): In any word, if vv and uu are two squares at position i and ww is a square at position $i + 1$, then either $|w| = |u|$ or $|w| = |v|$ or $|w| \geq 2|v|$.
- Consequently, in any infinite word, if infinitely many distinct squares begin at position i and ww is a square beginning at position $i + 1$, then $|w| = |u|$ for some square uu occurring at position i , so there is an overlap at position i .

Thank you!