

# Abelian Repetitions and Related Topics

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# Abelian repetitions

Erdős 1961 **abelian square**: a word  $xx'$  such that  $x'$  is a permutation of  $x$  (like reappear)

Evdokimov 1968 abelian squares avoidable over 25 letters

Pleasants 1970 abelian squares avoidable over 5 letters

Justin 1972 abelian 5-powers avoidable over 2 letters

Dekking 1979 abelian 4-powers avoidable over 2 letters

abelian cubes avoidable over 3 letters

Keränen 1992 abelian squares avoidable over 4 letters

# Dekking's construction

- ▶ Define a map:

$$a \rightarrow aaab, \quad b \rightarrow abb.$$

- ▶ The limit of the sequence

$$a \rightarrow aaab \rightarrow aaabaaabaaababb \rightarrow \dots$$

contains no abelian 4-power.

# Avoiding patterns in the abelian sense

- ▶ Avoiding the **pattern**  $xyyx$  in the **abelian sense** means avoiding all words  $xyy'x'$  where  $x$  and  $x'$  (resp.  $y$  and  $y'$ ) are permutations of each other.
- ▶ What are the patterns that are avoidable in the abelian sense?

# The Zimin patterns

- ▶ The **Zimin patterns**:

$$Z_1 = x, \quad Z_2 = xyx, \quad Z_3 = xyxzxyx, \quad \dots$$

## Conjecture (Currie and Linek 2001)

A pattern  $p$  with  $n$  variables is avoidable in the abelian sense if and only if  $Z_n$  avoids the pattern  $p$  in the abelian sense.

# Avoiding long binary patterns

## Theorem (Currie and Visentin 2008)

Any pattern over  $\{x, y\}$  of length greater than 118 is avoidable in the abelian sense over a 2-letter alphabet.

# The proof technique

- ▶ Let  $w$  be the infinite word

$aaaabaaaabaaaabaaaababbbaaaaab \dots$

generated by the map  $a \rightarrow aaaab$  and  $b \rightarrow abbb$ .

- ▶ Define a **weight function**  $f : \{a, b\}^* \rightarrow \mathbb{Z}_{11}$  by  $f(a) = 3$ ,  $f(b) = 10$ , and  $f(xy) = f(x) + f(y)$  for words  $x, y$ .
- ▶ Observe:

$$\{f(U) : U \text{ a prefix of } w\} = \{0, 1, 2, 3, 6, 9\}.$$

# An example (due to J. Currie)

- ▶  $w$  contains the pattern  $xyxyxy$  in the abelian sense:

*aaaabaabaabaabaabbbbaaab...*

- ▶ Write  $w = UX_1Y_1X_2Y_2Y_3X_3Y_4w'$  (here  $U = \epsilon$ ).
- ▶  $f(X_1) = f(X_2) = f(X_3) = 9$
- ▶  $f(Y_1) = f(Y_2) = f(Y_3) = f(Y_4) = 5$



# Computing the prefix weights

Prefix weights:

$$f(U) = 0, \quad f(UX_1) = 9, \quad f(UX_1Y_1) = 3,$$

$$f(UX_1Y_1X_2) = 1, \quad f(UX_1Y_1X_2Y_2) = 6,$$

$$f(UX_1Y_1X_2Y_2Y_3) = 0, \quad f(X_1Y_1X_2Y_2Y_3X_3) = 9,$$

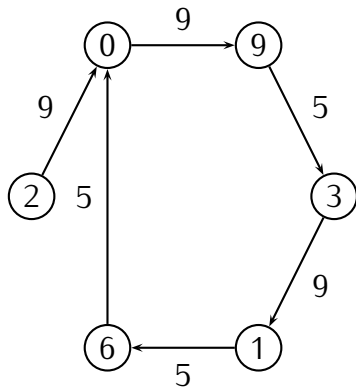
$$f(X_1Y_1X_2Y_2Y_3X_3Y_4) = 3$$

# A graph-theoretic viewpoint

View as a walk on the graph  $G(\{5, 9\})$  with

- ▶ vertex set  $\{0, 1, 2, 3, 6, 9\}$ , and
- ▶ edge set  $\{(i, j) : j - i \pmod{11} \in \{5, 9\}\}$ .

# The graph $G(\{5, 9\})$



# Concluding that long patterns are avoidable

- ▶ If  $w$  contains  $p$ , then  $p$  can be walked on some  $G(\{a, b\})$ .
- ▶ If  $p$  contains  $xxxx$ , it is avoidable.
- ▶ Show: all long patterns walkable on some  $G(\{a, b\})$  contain  $xxxx$ .
- ▶ Technical point: all long patterns should be walkable on  $G(\{a, b\})$  with  $a, b$  non-zero.

# Avoiding long abelian squares

## Theorem (Entringer, Jackson, Schatz 1974)

Any binary word of length  $k^2 + 6k$  contains an abelian square  $xx'$  with  $|x| \geq k$ .

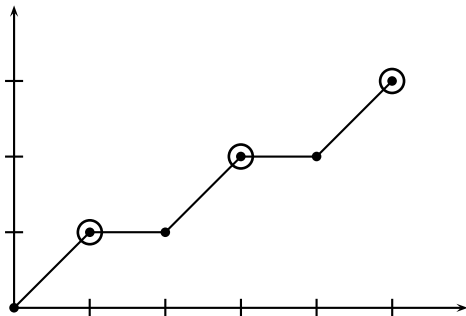
# A lower bound

## Theorem (Grant 2008)

Let  $q$  be a positive integer. There exists a binary word of length  $q(q + 1)$  which contains no abelian squares  $xx'$  with  $|x| \geq \sqrt{2q(q + 1)}$ .

# Geometrical interpretation

- ▶ Let  $w$  be a word. Let  $S_i = \sum_{j=1}^i w[j]$ . Plot  $(i, S_i)$ .
- ▶ abelian square  $\Rightarrow$  3 equidistant collinear points.
- ▶ For  $w = 10101$ :



# The construction

- ▶ Construct a sequence whose graph approximates a quadratic function.
- ▶ For  $0 \leq i \leq q(q+1)$ , define

$$a_i = \left\lfloor \frac{i^2}{2q(q+1)} \right\rfloor.$$

- ▶ Note:  $a_i - a_{i-1} \in \{0, 1\}$  for  $1 \leq i \leq q(q+1)$ .
- ▶ Define  $w$  by  $w_i = a_i - a_{i-1}$  for  $1 \leq i \leq q(q+1)$ .



# An example for $q = 3$

$$\begin{aligned}(a_i) &= ([1/24], [4/24], [9/24], \dots, [144/24]) \\ &= (0, 0, 0, 0, 0, 1, 1, 2, 2, 3, 4, 5, 6)\end{aligned}$$

$$(a_i - a_{i-1}) = (0, 0, 0, 0, 1, 0, 1, 0, 1, 1, 1, 1)$$

$$w = 000010101111$$

# Avoiding long abelian squares and cubes

## Problem (Mäkelä 2002)

Is there a constant  $c$  such that the set of abelian cubes  $xx'x''$  with  $|x| \geq c$  is avoidable over a 2-letter alphabet?

Is there a constant  $d$  such that the set of abelian squares  $xx'$  with  $|x| \geq d$  is avoidable over a 3-letter alphabet?

# A variation

## Theorem (Gerver and Ramsay 1979)

Let  $S$  be any set of non-collinear vectors in  $\mathbb{Z}^2$ . Any infinite  $S$ -walk contains arbitrarily large sets of collinear points.

# A variation

## Theorem (Gerver and Ramsay 1979)

Let  $S$  be any set of non-coplanar vectors in  $\mathbb{Z}^3$ . There is an infinite  $S$ -walk for which no  $5^{11} + 1$  points are collinear.

Can this be improved?

# Abelian complexity

- ▶ `roasting` and `organist` are **abelian equivalent** because they are anagrams.
- ▶ Let  $w$  be an infinite word.
- ▶ Define  $\rho^{\text{ab}}(n) =$  the number of abelian equivalence classes of words of length  $n$  in  $w$ .
- ▶ This is the **abelian complexity function** of  $w$ .

# The binary case

## Theorem (Coven and Hedlund 1973)

A recurrent infinite word has abelian complexity  $\rho^{\text{ab}}(n) = 2$  for all  $n \geq 1$  if and only if it is Sturmian.

# The ternary case

## Problem (Rauzy 1982/83)

Does there exist a recurrent infinite word with abelian complexity  $\rho^{\text{ab}}(n) = 3$  for all  $n \geq 1$ ?

Richomme, Saari, and Zamboni answered Rauzy's question in the affirmative.

# The 4-letter case

Problem (Richomme, Saari, Zamboni 2009)

Does there exist a recurrent infinite word with abelian complexity  $\rho^{\text{ab}}(n) = 4$  for all  $n \geq 1$ ?



# Larger alphabets

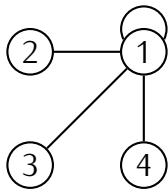
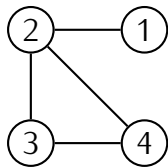
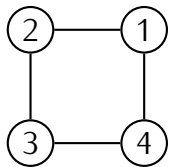
## Theorem (Currie and R. 2009)

Let  $k \geq 4$ . There is no recurrent infinite word with abelian complexity  $\rho^{\text{ab}}(n) = k$  for all  $n \geq 1$ .

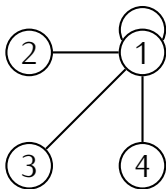
# The idea of the proof

- ▶ Suppose  $w$  is a recurrent word over a 4-letter alphabet with constant abelian complexity 4.
- ▶ Define a graph  $G$  with:
  - ▶ vertex set  $\{1, 2, 3, 4\}$ , and
  - ▶ edges  $ij$  if either  $ij$  or  $ji$  is a factor of  $w$ .
- ▶  $G$  is a spanning tree with one extra edge.

# Examples of the graph $G$



# A particular case



- ▶ Suppose 111 a factor of  $w$ .
- ▶ Note: 213, 214, 312, etc. not factors.
- ▶ Let  $m$  be minimal such that  $d1^m e$  is a factor (wlog  $21^m 3$  is a factor).

# A particular case

- ▶ Then

$$121^m 3, 1^m 21^2, 1^m 31^2, 1^m 41^2$$

are the 4 factors of length  $m + 3$  (up to permutation).

- ▶ Note:  $1^{m+2}$  not a factor.
- ▶ Consider a factor  $b1^k c$  where  $b \in \{2, 3\}$  and  $c = 4$ .
- ▶ Either  $k = m$  or  $k = m + 1$ . Both cases give a 5th factor. Contradiction.
- ▶ There are many other cases.

# Conclusion

- ▶ Many questions already resolved for ordinary repetitions remain open for abelian repetitions.
- ▶ Problems regarding abelian repetitions are much more difficult.
- ▶ Abelian complexity: a relatively recent concept. Probably much more can be done in this area.

The End